

# Bicategories of algebras for relative pseudomonads

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## 1 Relative (pseudo)monads & the presheaf construction

**Definition 1.1** (Relative monad [1]). Let  $J: \mathbb{C} \rightarrow \mathbb{D}$  be a functor between categories. A  $J$ -relative monad  $(T, i, (-)^*)$  comprises

- *units*  $i_X: JX \rightarrow TX$  in  $\mathbb{D}$  for  $X \in \text{ob } \mathbb{C}$ ;
- *extensions*  $f^*: TX \rightarrow TY$  in  $\mathbb{D}$  for  $f: JX \rightarrow TY$ ,

satisfying  $f = f^*i$ ,  $(f^*g)^* = f^*g^*$ ,  $i_X^* = 1_{TX}$  for  $f: JX \rightarrow TY$  and  $g: JW \rightarrow TX$ .

If  $J$  is instead a pseudofunctor between *bicategories*, we may weaken the above equalities to invertible 2-cells, obtaining a *relative pseudomonad* [2].

**Example 1.2.** (Presheaves) The presheaf construction  $P: \text{Cat} \rightarrow \text{CAT}$  has the structure of a relative pseudomonad. The units are Yoneda embeddings  $y_X: X \rightarrow PX$ , and the extension of  $f: X \rightarrow PY$  is given by its left extension along  $y_X$  (i.e. by taking a certain colimit):

$$\begin{array}{ccc} X & \xrightarrow{y_X} & PX \\ & \searrow f \cong & \downarrow f^* := y_X \triangleright f \\ & & PY \end{array}$$

## 2 The bicategory of pseudoalgebras for a relative pseudomonad

The Kleisli bicategory for a relative pseudomonad was constructed in [2]; we construct the bicategory of pseudoalgebras.

**Definition 2.1** ( $T$ -pseudoalgebra). For a  $J$ -relative pseudomonad  $T$ , a  $T$ -pseudoalgebra  $(A, (-)^a; \tilde{a}, \hat{a})$  comprises

- an *underlying object*  $A$  in  $\text{ob } \mathbb{D}$ ;
- *algebra extensions*  $f^a: TX \rightarrow A$  in  $\mathbb{D}$  for  $f: JX \rightarrow A$ ,

along with two families of invertible 2-cells  $\tilde{a}_f: f \rightarrow f^a i_X$  and  $\hat{a}: (f^a g)^a \rightarrow f^a g^a$  for  $f: JX \rightarrow A$  and  $g: JW \rightarrow TX$ .

When  $J$  is the identity, this definition reduces to that of a pseudoalgebra for a *no-iteration pseudomonad* as defined in [3].

Defining notions of algebra morphism and algebra transformation, we may assemble pseudoalgebras into a bicategory.

**Theorem 2.2.** Let  $T$  be a relative pseudomonad.  $T$ -pseudoalgebras, pseudomorphisms and transformations form a bicategory  $\text{PsAlg}(T)$ , called the bicategory of  $T$ -pseudoalgebras.

To demonstrate that our construction is correct, we find that  $\text{PsAlg}(T)$  satisfies the following universal property.

**Theorem 2.3.** Let  $T$  be a  $J$ -relative pseudomonad. The bicategory of pseudoalgebras and pseudomorphisms forms a relative pseudoadjunction which is bi-terminal among resolutions of  $T$  (the  $J$ -relative pseudoadjunctions that induce  $T$ ).

$$\begin{array}{ccc} & \text{PsAlg}(T) & \\ F_T \nearrow & \dashv & \searrow U_T \\ \mathbb{C} & \xrightarrow{J} & \mathbb{D} \end{array}$$

Likewise, the Kleisli resolution is bi-initial, and there is an canonical fully faithful pseudofunctor  $I_T: \text{Kl}(T) \rightarrow \text{PsAlg}(T)$  which is the embedding of the free  $T$ -pseudoalgebras.

## 3 Idempotence does not imply algebraic idempotence

We explore two notions of idempotence for relative monads.

**Definition 3.1.** (Idempotent relative monad) A relative monad  $(T, i, (-)^*)$  is *idempotent* if the extension operator  $(-)^*: \mathbb{D}(JX, TY) \rightarrow \mathbb{D}(TX, TY)$  is inverse to precomposition with the unit  $- \circ i_X: \mathbb{D}(JX, TY) \rightarrow \mathbb{D}(TX, TY)$ , for all  $X, Y$  in  $\text{ob } \mathbb{C}$ .

**Definition 3.2.** (Algebraically idempotent relative monad) A relative monad  $(T, i, (-)^*)$  is *algebraically idempotent* if the algebra extension  $(-)^a: \mathbb{D}(JX, A) \rightarrow \mathbb{D}(TX, A)$  is inverse to precomposition with the unit  $- \circ i_X: \mathbb{D}(JX, A) \rightarrow \mathbb{D}(TX, A)$ , for all  $X$  in  $\text{ob } \mathbb{C}$  and  $(A, (-)^a)$  in  $\text{Alg}(T)$ .

Algebraic idempotence implies idempotence (which, by definition, is algebraic idempotence restricted to free algebras). In fact, in the special case where  $J$  is an identity, the two notions coincide. However, in general they do not. A counterexample is given by the following small category  $\mathbb{D}$ ,

$$J \xrightarrow{i} T \begin{array}{c} \xrightarrow{h} \\ \xleftarrow{g} \end{array} A$$

subject to  $gi = hi$ . Considering  $J$  and  $T$  as functors from the point to  $\mathbb{D}$ ,  $T$  is a  $J$ -relative monad which is idempotent but not algebraically idempotent.

## 4 The presheaf construction is algebraically lax-idempotent

In the bicategorical setting, we may generalise from (algebraic) idempotence to (*algebraic*) *lax-idempotence*, asking only for an adjunction  $(-)^* \dashv - \circ i_X$  or  $(-)^a \dashv - \circ i_X$ . The presheaf relative pseudomonad  $P$  is lax-idempotent by construction; however, showing it to be algebraically lax-idempotent is (perhaps surprisingly) nontrivial.

**Proposition 4.1.** Let  $(A, (-)^a)$  be a  $P$ -pseudoalgebra. Then  $A \in \text{CAT}$  is a cocomplete category.

The colimit of a diagram  $f: D \rightarrow A$  can be computed as the image of the terminal presheaf under  $f^a: PD \rightarrow A$ .

We can also show that every pseudomorphism of algebras  $(A, (-)^a) \rightarrow (B, (-)^b)$  is cocontinuous as a functor  $A \rightarrow B$ . Consequently, we have the following.

**Theorem 4.2.** The presheaf relative pseudomonad  $P$  is algebraically lax-idempotent.

**Corollary 4.3.** The bicategory  $\text{Dist}$  of small categories and distributors is biequivalent to the full sub-2-category of the 2-category of locally small cocomplete categories and cocontinuous functors, spanned by the presheaf categories.

## References

- [1] Thomas Altenkirch, James Chapman, and Tarmo Uustalu. Monads need not be endofunctors. *Logical Methods in Computer Science*, 11(1), 2015.
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- [3] Francisco Marmolejo and Richard J Wood. No-iteration pseudomonads. *Theory and Applications of Categories*, 28(14):371–402, 2013.